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1	TITLE OF THE INVENTION
2	Minimum Cost Routing Based on Relative Costs of Node Resources
3	BACKGROUND OF THE INVENTION
4	Field of the Invention
5	The present invention relates generally to communications networks
6	where network nodes are interconnected by a number of links, and more
7	specifically to a least-cost routing technique. The present invention is
8	particularly useful for optical networks.
9	Description of the Related Art
10	In a prior art communications network as shown and described in
11	Japanese Patent Publications 07-250356 and 10-276208, neighboring nodes are
12	interconnected by a plurality of communication links that terminate a
13	switching system of the nodes. Each node maintains a database that stores
14	the status information of the links. A path controller exchanges the link
15	status information with the neighboring nodes to update their database and
16	determines a minimum cost route in response to a path setup request from a
17	client device to produce a path setup signal. The switching system responds
18	to the path setup signal for establishing a connection between the links
19	specified by the determined route.
20	However, since the prior art is solely based on the costs associated
21	with communication links, the route determined is not always the best for
22	nodes involved in a given path if their node resources are different in
23	quantity, size and type.
24	SUMMARY OF THE INVENTION
25	It is therefore an object of the present invention to improve the prior

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art routing decision by taking into account the per-channel costs of node resources.

According to a first aspect of the present invention, there is provided a network node for a communications network which comprises a plurality of network nodes interconnected by communication links, wherein the network node is one of the plurality of network nodes. The network node comprises a first module having a plurality of input ports and a plurality of output ports for handling a group of channels between the input ports and the output ports as a routing unit and a second module having an input port and an output port for handling a channel between the input port and the output port as the routing unit. A module state and link state databases are provided. In the module state database, module cost data of the first and second modules and module cost data of other network nodes are stored in the link state database link cost data of the communication links. Further provided is switching system for determining a route of minimum cost by using the module state database and the link state database and establishing, according to the determined route, a connection between one of a plurality of incoming communication links and one of the input ports of the first and second modules and establishing a connection between one of the output ports of the first and second modules and one of a plurality of outgoing communication links. According to a second aspect, the present invention provides a communications network comprising a plurality of network nodes interconnected by incoming links and outgoing links. Each of the network nodes comprises a first module for handling a group of channels between its

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input ports and its output ports as a routing unit, a second module for 1 handling a channel between its input port and its output port as a routing 2 unit. A module state database stores module cost data of the first and second 3 modules and module cost data of other nodes of the network and a link state 4 database stores link cost data of the incoming and outgoing links. A 5 switching system determines a route of minimum cost by using the module 6 state database and the link state database and establishes, according to the 7 determined route, a connection between one of said incoming links and one 8 of the input ports of the first and second modules and establish a connection 9 between one of the output ports of the first and second modules and one of 10 said outgoing links. Terminating circuitry is provided for transmitting an 11 advertisement message to neighboring network nodes for communicating the 12 contents of the module state database and receiving an advertisement 13 message from the neighboring network nodes for updating the module state 14 database according to the received message. 15 16

According to a third aspect, the present invention provides a centralized network management system comprising a plurality of network nodes interconnected by a plurality of incoming communication links and a plurality of outgoing communication links, a management center connected to the network nodes via control channels. Each of the network nodes comprises a first module for handling a group of channels between its input ports and its output ports as a routing unit, a second module for handling a channel between its input port and its output port as a routing unit. A switching fabric having a plurality of incoming interfaces for interfacing incoming communication links and a plurality of outgoing interfaces for

to the following drawings, in which:

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1 interfacing outgoing communication links. The management center 2 comprises a module state database, a link state database and a path 3 controller. The path controller store module cost data of the first and second modules of each of the network nodes into the module state database and 4 5 link cost data of the incoming and outgoing communication links into the 6 link state database and determines a route of minimum cost by using the 7 module and link state databases. Path controller controls the switching fabric 8 of each network node to establish a connection between one of the incoming 9 interfaces and one of the input ports of the first and second modules and a 10 connection between one of the output ports of the first and second modules 11 and one of the outgoing interfaces. 12 According to a fourth aspect, the present invention provides a routing 13 method comprising the steps of providing, in each of a plurality of network 14 nodes interconnected by communication links, a first module having a 15 plurality of input ports and a plurality of output ports, the first module 16 handling a group of channels between the input ports and the output ports as 17 a routing unit, providing a second module having an input port and an 18 output port, the second module handling a channel between the input port 19 and the output port as the routing unit, storing module cost data of the first 20 and second modules in a module state database, storing link cost data of the 21 communication links in a link state database, and determining a route of 22 minimum cost by using the module state database and the link state 23 database. 24 BRIEF DESCRIPTION OF THE DRAWIGNS 25 The present invention will be described in detail further with reference

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1	Fig. 1 is a block diagram of an optical communications network
2	according to a first embodiment of the present invention;
3	Fig. 2 is an illustration of the module state database of Fig. 1;
4	Fig. 3 is an illustration of the link state database of Fig. 1;
5	Fig. 4 is a flowchart of the operation of a path controller when
6	transmitting a link or a module state advertisement message;
7	Fig. 5 is a flowchart of the operation of the path controller when
8	receiving the link or module state advertisement message;
9	Fig. 6 is a flowchart of the operation of the path controller when a
10	route of minimum cost is determined and a path is set up according to the
11	determined route;
12	Fig. 7 is a is a block diagram of an optical communications network
13	according to a second embodiment of the present invention;
14	Fig. 8 is an illustration of the module state database of Fig. 7;
15	Fig. 9 is an illustration of the link state database of Fig. 7; and
16	Fig. 10 is a block diagram of an optical communications network
17	according to a third embodiment of the present invention.
18	DETAILED DESCRIPTION
19	An optical network according to a first embodiment of the present
20	invention is shown in Fig. 1. The network comprises a plurality of optical
21	nodes 1, 2 and 3 interconnected by optical links for transporting traffic signals
22	and control channels for transmission of control messages. As a
23	representative node, details of node 2 are illustrated.
24	According to the definition of the present invention, a "path" is a
25	transmission unit established between two nodes via at least one

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1 intermediate node, a "link" is a transmission medium established between 2 two neighboring nodes, a "forwarding adjacency link" is a virtual 3 transmission medium established between two nodes via at least one intermediate node, and a "channel" is a logical transmission unit established 4 5 between two neighboring nodes. By using wavelength multiplexing 6 technique, a plurality of channels can be accommodated in a single link and a 7 plurality of paths can be accommodated in a single forwarding adjacency 8 link. 9 Node 2 is comprised of a switch fabric 25 and a number of node resources implemented with manufacturing units, or modules such as a 10 group switch module (GSM) 26, a channel switch module (CSM) 27 and a 11 12 group regenerator module (GRM) 28. These modules are classified as group-13 associated modules and channel-associated modules. The group-associated 14 modules are designed to handle a group of channels as a routing unit, while the channel-associated modules are designed to handle channels as 15 16 individual routing units. The group-associated modules are implemented 17 with as low a production cost as possible for a given volume of traffic, so that 18 their per-channel cost is less than the corresponding cost of the channel-19 associated modules. Therefore, it is advantageous to use the groupassociated modules for applications in which a heavy traffic volume can be 20 21 divided into a group of channels of same route, whereas the channel-22 associated modules may be advantageous for low-traffic applications. 23 Therefore, the group switch module 26 has a lower per-channel cost 24 than the corresponding cost of the channel switch module 27. Group

regenerator module 28 is a product that is manufactured as a group-

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1	associated module and includes, for example, two 3R (reshaping, retiming,
2	regeneration) regenerators. Thus, the per-channel cost of GRM 28 is less than
3	the per-channel cost of an individual 3R regenerator.
4	As will be described, the present invention provides least cost route
5	decision based on the per-channel cost of node resources.
6	Node 1 has outgoing (OG) interfaces OG10 ~ OG17 which are
7	connected to corresponding incoming (IC) interfaces IC20 ~ IC27 of node 2
8	using optical links $10 \sim 17$. Eight optical links $30 \sim 37$ connect outgoing
9	interfaces OG20 ~ OG27 of node 2 to incoming interfaces IC30 ~ IC37 of node
10	3. Each optical link carries a single wavelength channel which is one of four
11	wavelengths $\lambda 1$, $\lambda 2$, $\lambda 3$ and $\lambda 4$, as illustrated. A client device 4 is provided,
12	which is terminated at an incoming interface IC28 and an outgoing interface
13	OG28 of node 2.
14	Switched connections are established in the switching fabric 25 so that
15	the incoming interfaces IC20 ~ IC28 can be connected to any of the input
16	ports of node resources 26, 27 and 28 and the output ports of these node
17	resources can be connected to any of the outgoing interfaces OG20 ~ OG28.
18	However, outgoing interfaces of the switching fabric 25 are selected
19	depending on whether the node resources have the ability to provide
20	wavelength conversion and on which incoming interfaces are used.
21	In the group switch module 26, the input (IP) ports IP1 ~ IP4 are
22	divided into two input port groups IPG1 and IPG2 and the output (OP) ports
23	of OP1 ~ OP4 are likewise divided into two output port groups OPG1 and
24	OPG2. In the illustrated example, incoming group connections 40 are
25	established in the switching fabric 25 for connecting the incoming interfaces

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1	IC20, IC21 to the input ports IP1, IP2 of the group switch module 26 and
2	outgoing group connections 42 for connecting the output ports OP1, OP2 of
3	the module 26 to the output interfaces OG22, OP23. Simultaneously with the
4	establishment of connections 40 and 42, the group switch module 26 sets up
5	connections 41 between the input ports IP1, IP2 and the output ports OP1,
6	OP2. Note that the connection P6 indicated by dotted line 40 is a standby
7	path which can be used as what is called a "forward adjacency link" which
8	remains in the established state for future demand from users. Accordingly,
9	when at least one path is established through the group switch module 26,
10	two incoming connections and two outgoing connections are simultaneously
11	established in the switching fabric 25 no matter whether all of the established
12	connections are immediately used for carrying traffic.
13	Channel switch module 27 establishes a channel connection between
14	one of its input ports IP5 ~ IP8 and one of its output ports OP5 ~ OP8. In one
15	example, connections 43, 46 and 45, 48 are established in the switching fabric
16	25 for connecting the incoming interfaces IC22, IC24 to the input ports IP5,
17	IP7 and connecting the output ports OP5, OP7 to the output interfaces OG24,
18	OP26. In this case, the channel switch module 27 may establish channel
19	connections 44 and 47 between the input ports IP5, IP7 and the output ports
20	OP5, OP7. It is assumed that the channel switch module 27 has the ability to
21	convert the wavelength of any of its channels at the input ports IP5 to IP8 to
22	the wavelength of desired outgoing interfaces OG.
23	In the case of the group regenerator module 28, group connections 49
24	and 50 are established in the switching fabric 25 to connect the incoming
25	interfaces 26, 27 to the input ports IP9, IP10 of group IPG3 of the module 28

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and connect the output ports OP9, OP10 of group OPG3 to the outgoing 1 2 interfaces OG20, OG21. Switching fabric 25 is controlled by a path controller 21. Preferably, 3 the group and channel witch modules 26 and 27 are also controlled by the 4 5 path controller 21. Node 2 maintains a module (or node) state database 22 6 and a link state database 23. As will be described in detail, the module state database 22 is controlled by the path controller 21 to store module status data 7 8 indicating the status of resources of nodes 1, 2 and 3 when a new module is 9 installed and update the stored module status data when an existing module 10 is removed or modified. Path controller 21 transmits a module state 11 advertisement message to the network to inform the neighboring nodes of the 12 new status of the node modules. 13 In like manner, the link state database 23 is controlled by the path 14 controller 21 to store link state data of optical links of nodes 1, 2 and 3 when a 15 new link is installed as well as link state data of a link known as a 16 "forwarding adjacency link" which connects nodes 1 and 3. When an 17 existing link is removed or modified, the database 23 is controlled to update 18 the stored link state data. Path controller 21 transmits a link state 19 advertisement message to the network to inform neighboring nodes of the 20 new status of the links. Therefore, the link and module status information 21 maintained by the nodes 1, 2 and 3 are compared with each other to share the 22 same knowledge of network resources. Path controller 21 of node 2 exchanges link and module state 23 24 advertisement messages with nodes 1 and 3 over control channels established 25 by a terminating unit 24. Terminating unit 24 also receives a request message

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1	from the client device 4. The received request message is processed by the
2	path controller 21 to determine and establish a path through the network
3	according to the link and module state database.
4	Fig. 2 shows details of the module state database 22. Module state
5	database 22 is divided into a plurality of entries corresponding to different
6	node modules of its own node as well as the node modules of other nodes of
7	the network. Each entry of the module state database 22 is divided into a
8	plurality fields for indicating attributes of the module (such as module
9	number, module type and wavelength convertibility) and a path number
10	field for indicating path numbers of the module. The entry of each module is
11	further divided into a plurality of input-port fields and a plurality of output-
12	port fields for indicating the various attributes of the input and output ports,
13	such as port number, port group number, port cost, available wavelength,
14	available data format, maximum data rate, and interface number. Each of the
15	input- and output-port fields is sub-divided into a number of sub-fields
16	corresponding to path numbers for specifying attributes of the input and
17	output ports associated with the path numbers.
18	The following is a description of the contents of the module state
19	database 22 of Fig. 2 by assuming that switched connections are set up in the
20	switching fabric 25 as shown in Fig. 1.
21	Entry 1 of module state database 22 contains the information of group
22	switch module 26, including its module number, module type and
23	unavailability of wavelength conversion. The input-port fields of entry $1\ \mathrm{are}$
24	divided into four sub-fields corresponding to input ports IP1, IP2, PI3 and
25	IP4.

In the entry 1, the sub-fields of input port IP1 indicate that port IP1

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belongs to port group IPG1 and that a path identified by path number P1 is 1 established from the incoming interface IC20 and the normalized cost of this 2 input port is "1" and the available wavelength, the available data format and 3 the maximum data rate of the incoming wavelength channel at input port IP1 4 can be arbitrarily determined. The sub-fields of output port OP1 indicate that 5 port OP1 belongs to port group OPG1 and that a path P1 is established from 6 this output port to the outgoing interface OG24 and the normalized cost of 7 8 this input port is "1" . The available wavelength, the available data format 9 and the maximum data rate of the incoming wavelength channel at the output port OP1 determined dependent on the values of the corresponding 10 11 input ports. For a path P6 established on a forwarding adjacency link 12 between nodes 1 and 3, the attributes of the group switch module 26 are 13 indicated by specifying the input port number IP2, the incoming interface 14 number IC22, the port cost =1, and the output port number OP2, the outgoing 15 interface number OG22 and the port cost = 1. 16 In the sub-fields of input port IP2 which also belongs to port group 17 IPG1, the corresponding path number sub-field is left blank to indicate that 18 no path is established for IP2, but a connection is established in the switching 19 fabric 25 from the incoming interface IC22 to input port IP2. The normalized 20 cost of input port IP is "1" and the available wavelength, the available data format and the maximum data rate of the incoming wavelength channel at 21 22 input port IP2 are left blank. In the sub-fields of output port OP2 which 23 belongs to port group OPG1, the corresponding path number sub-field is left 24 blank to indicate that no path is established for OP2, but a connection is established in the switching fabric 25 from output port OP2 to the outgoing 25

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interface OG22. The normalized cost of input port OP is "1" and the available 1 2 wavelength, the available data format and the maximum data rate of the 3 incoming wavelength channel at input port IP2 are left blank. 4 Entry 2 of database 23 contains the information of channel switch module 27, including its module number, module type and availability of 5 6 wavelength conversion. The input-port fields of entry 2 are divided into four 7 sub-fields corresponding to input ports IP5, IP6, PI7 and IP8. The input port 8 cost of the channel switch module 27 is "2", for example, and the available 9 wavelength, data format and maximum data rate of the channel switch 10 module 27 are predetermined. For input port IP5, the associated sub-fields 11 contain information indicating that a path P2 is established from the 12 incoming interface IC21 to the input port IP5. Similarly, the sub-fields of 13 input port IP6 contain information indicating that a path P3 is established 14 from the incoming interface IC24 to the input port IP6. The output port cost 15 of the channel switch module 27 is "2", for example, and the available data 16 format and maximum data rate of the channel switch module 27 are 17 predetermined. However, the available wavelength can be arbitrarily 18 determined because of the wavelength conversion capability of the module 19 27. The sub-fields of output port OP5 contain information indicating that the 20 path P2 is established from output port OP5 to the outgoing interface OG23. 21 Similarly, the sub-fields of output port OP6 contain information indicating 22 that the path P3 is established from port OP6 to the outgoing interface OG20. 23 Entry 3 of database 23 contains the information of group regenerator 24 module 28, including its module number, module type and availability of 25 wavelength conversion. The input-port fields of entry 3 are divided into two

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1 sub-fields corresponding to input ports IP9 and IP10. The normalized input 2 port cost of the group regenerator module 28 is "1", and the available 3 wavelength, data format and maximum data rate of the module 28 are predetermined. Sub-fields of input port IP9 indicate that a path P4 is 4 5 established from the incoming interface IC26 to the input port IP9. Similarly, 6 the sub-fields of input port IP10 indicate that a path P5 is established from the 7 incoming interface IC27 to the input port IP10. The output port cost of the group regenerator module 28 is "1", and the available wavelength, data 8 9 format and maximum data rate of the module 28 are predetermined. The sub-fields of output port OP9 indicate that the path P4 is established from the 10 11 output port OP9 to the outgoing interface OG26. Similarly, the sub-fields of 12 output port OP10 indicate that the path P5 is established from the port OP10 13 to the outgoing interface OG27. 14 Details of the link state database 23 are shown in Fig. 3. Link state 15 database 23 of node 2 is divided into a number of link entries corresponding 16 to links between nodes 1 and 2 and links between nodes 2 and 3 and 17 forwarding adjacency links between nodes 1 and 3 via node 2. Each of the 18 link entries is sub-divided into a plurality of fields to define a number of 19 attributes of each link. The defined attributes of a link include the node 20 addresses of transmit and receive nodes, the outgoing and incoming interface 21 numbers, the link cost, the wavelength of the link, the data format, the 22 maximum data rate and the path number. 23 The link entry 1 of database 23, for example, specifies attribute data of 24 link 10 by giving the addresses of nodes 1 and 2, outgoing (OG) interface 25 number = OG10, incoming (IC) interface number = IC20, link cost = 1,

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wavelength = $\lambda 1$, and path number = P1. For path P6 established on a 1 2 forwarding adjacency link between nodes 1 and 3, a link entry 2 is defined to 3 specify the attributes of the forwarding adjacency link by indicating the addresses of transmit and receive nodes 1 and 3, outgoing interface number = 4 5 OG12, incoming interface number = IC32, link cost = 4 (=1 + 2 + 1), wavelength = $\lambda 1$, and path number = P6. The link cost includes the total cost 6 7 of links 10 and 30 plus the total cost of input port IP1 and output port OP1 of 8 group switch module 26. 9 Since the forwarding adjacency link entries include data of total 10 forwarding adjacency link cost, route computations can be simplified. 11 The present invention is further characterized by the transmission of 12 the module and link state advertisement messages from the path controller to 13 neighboring nodes and the reception of these messages for creating or 14 updating an entry in the databases of the neighboring nodes. The link and 15 module state advertisement messages are hereinafter called "resource" state 16 advertisement messages since the process associated with transmission and 17 reception of a module state advertisement message is identical to the process 18 associated with transmission and reception of a link status advertisement 19 message. Further, links and node modules are collectively called "resources" 20 and the link and module state databases are called "resource state databases". 21 Fig. 4 is a flowchart of the operation of path controller 21 of node 2 22 when transmitting a resource state advertisement message. When a new 23 resource is installed or established in the node 2, the decision at step 401 is 24 affirmative and flow proceeds to step 402 to add a new entry to the 25 corresponding resource database for storing status data associated with the

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new resource. If the decision at step 401 is negative or a new entry is created 1 2 at step 402, flow proceeds to decision step 403 to check to see if an existing 3 resource is removed. If this is the case, an entry corresponding to the removed resource is removed from the corresponding resource state 4 database. If the decision at step 403 is negative or an entry is removed at step 5 404, flow proceeds to decision step 405 to check to see if the status data of an 6 7 existing resource is modified. If so, a corresponding entry of the database is 8 modified at step 406. 9 The next decision is to determine if a resource state advertisement 10 message is to be sent to the network at step 407. If the path controller has 11 executed step 402, 404 or 406, the controller proceeds to step 408 and 12 formulates a resource state advertisement message with necessary data and 13 sends it to the nodes 1 and 3 via the terminating unit 24. 14 In Fig. 5, the path controller of each node monitors the associated 15 terminating unit 24 to receive a resource state advertisement message. When 16 such a message is received (step 501), the controller 21 determines whether it 17 contains a new resource entry, a removed entry or a modified entry by 18 comparing the contents of the advertisement message with those of the 19 corresponding database. If new resource entry data is contained in the 20 received advertisement message (step 502), a new entry is created in the corresponding resource state database (step 503). If the comparison indicates 21 22 that an existing entry of the database is missing in the received advertisement 23 message, the existing entry is removed from the database (step 505). If the 24 comparison indicates that differences exist between an entry of the 25 advertisement message (step 506) and the corresponding entry of the

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database, the latter is modified according to the entry of the message (step 1 2 507). At step 508, the path controller 21 retransmits the received 3 advertisement message to adjacent nodes, and returns to the starting point of the routine. 4 Since the data concerning the attributes of resources (such as module 5 number, module type, port number, port group number, wavelength, data 6 format, data rate) remains relatively static, while the data concerning the 7 8 attributes of paths vary from instant to instant, the data concerning the 9 attributes of resources are formulated in a resource (module/link) state 10 advertisement message and flooded to all nodes of the network. On the other 11 hand, the data concerning the path attributes are formulated in a signaling 12 message and transmitted to downstream nodes. The processes associated 13 with the transmission and reception of the signaling message are identical to 14 those described above with reference to the resource state advertisement 15 message. 16 When the client device 4 establishes a path through the network, it 17 sends a path setup message to the path controller 21. The process performed 18 by the path controller 21 during a path setup mode is shown in Fig. 6. When 19 the path controller receives a path setup message (step 601), it searches the 20 link state database 23 for links that are available along possible routes to the 21 destination node specified in the setup message (step 602). If available routes 22 are found (step 603), the path controller makes a search through the module 23 state database 22 for available node modules (step 604). If available node 24 modules are found (step 605), the path controller proceeds to step 606 to form 25 candidate paths using the available links and the available node modules. A

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total cost of each of the candidate paths is calculated (step 607) and all total 1 2 costs are compared and one of the candidate paths having a minimum value is determined as a path to the destination (step 608). At step 609, connections 3 are established in the switching fabric 25 according to the routing data of the 4 determined path. Path controller 21 formulates with the routing data of the 5 path and transmits a path setup message to a downstream node to cause it to 6 establish connections (step 610). If the downstream node is an intermediate 7 node of the path, it will reformulate a path setup message and transmits it 8 9 downstream. 10 In the previous embodiment, databases 22 and 23 are managed on a 11 per-channel basis, and route calculation and decision are performed on a perchannel basis. 12 13 According to a second embodiment of the present invention, channels 14 of different wavelengths are multiplexed onto inter-node optical links (physical links). A series of concatenated links, or forwarding adjacency (FA) 15 16 links (virtual links) are established between two network nodes via at least 17 one intermediate node. Each virtual link can transport at least one 18 wavelength channel. With respect to the physical and virtual links, a group 19 of wavelength channels is handled as a unit for making a routing decision. In 20 each of the network nodes the group switch module is a link switch designed 21 to handle a group of wavelength channels as a route decision unit, and the 22 channel switch module is designed to handle each wavelength as a route 23 decision unit. Note that the input port of the channel switch module has the 24 ability to tune to any of the multiplexed incoming wavelengths and the 25 output port has the ability to tune to any of the multiplexed outgoing

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1	wavelengths.
2	One example of the second embodiment of this invention is shown in
3	Fig. 7. Nodes 1 and 2 are connected by optical links 10 ~ 13 and nodes 2 and
4	3 are connected by optical links 30 ~ 33. A pair of wavelengths $\lambda 1$ and $\lambda 2$ or a
5	pair of wavelengths $\lambda 3$ and $\lambda 4$ are multiplexed onto each optical link, as
6	illustrated. Node 2 includes a link switch module 51, a channel switch
7	module 52 and a channel-associated regenerator 53.
8	Link switch module 51 has two input ports IP1 and IP2 and two
9	output ports OP1 and OP2 for connecting any of the input ports to any of the
10	output ports on a per-link basis, i.e., on a per-multiplex group basis.
11	Channel switch module 52 has two input ports IP3 and IP4 and two
12	output ports OP3 and OP4 for connecting any of the input ports to any of the
13	output ports on a per-channel basis. Each of the input ports IP3 and IP4 has
14	the function of tuning to any of the incoming wavelengths $\lambda 1\sim \lambda 4$ and each
15	of the output ports OP3 and OP4 has the function of tuning to any of the
16	outgoing wavelengths $\lambda 1 \sim \lambda 4$.
17	A 3R regenerator 53 is provided as a channel-associated module
18	having an input port IP5 and an output port OP5.
19	In response to a switching control signal from the path controller 21,
20	the switching fabric 25 establishes a connection between any of its incoming
21	interfaces IC20 ~ IC24 to any of the input ports IP1 ~ IP5 and a connection
22	between any of the output ports OP1 ~ OP5 to any of its outgoing interfaces
23	OG20 ~OG24. Since the link switch module 51 is capable of carrying four
24	wavelength channels, the per-channel cost of this module is lower than the
25	corresponding cost of the channel switch module 52.

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1 Path controller 21 formulates the switching control signal using 2 module and link status data stored in the databases 22 and 23. Details of the module state database 22 and link state database 23 of Fig. 7 are respectively 3 4 shown in Figs. 8 and 9 by assuming that the node 2 establishes a virtual link 5 60 of channels $\lambda 1$ and $\lambda 2$ between nodes 1 and 3 with connections 61, 62 and 63 and a virtual link 63 of channel $\lambda 3$ with connections 65, 66 and 67, and that 6 7 path numbers P1, P2 and P5 are assigned to channels λ 1, λ 2 and λ 3, 8 respectively. 9 As shown in Fig. 8, the module state database 22 of Fig. 7 defines 10 entries 1, 2 and 3 for modules 51, 52 and 53, respectively. Each entry lists 11 module attributes including module number, module type and wavelength 12 convertibility and for each module entry one or more path attributes are 13 defined, with the path attributes being divided into those associated with the 14 input port of the module and those associated with the output port. For entry 15 1 of the link switch module 51, path numbers P1 and P2 are specified for 16 channels $\lambda 1$ and $\lambda 2$. Port cost of the input port number IP1 is set equal to 1 17 and the port cost of the output port number OP1 is set equal to 1. In the entry 18 2, the status data associated with the channel switch module 52 indicates that 19 the path number is P5 and the channel wavelength is $\lambda 3$. Port cost of the 20 input port number IP3 is set to 2 which is higher than the port cost of module 21 51 and the port cost of the output port number OP3 is also set to 2. 22 In Fig. 9, the link state database 23 of Fig. 7 is divided into a plurality of physical link entries and a plurality of virtual link entries. The physical 23 24 link entries are associated individually with the incoming optical links 10 ~ 13 and the outgoing optical links 30 ~ 33. Each physical link entry defines in the 25

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1	node address fields the attributes of a physical link between adjacent nodes 1
2	and 2 or a physical link between adjacent nodes 2 and 3. The virtual link
3	entries are associated with virtual links established between the nodes 1 and
4	3 via node 2. Each virtual link includes multiplexed wavelength channels
5	and defines a virtual link between nodes 1 and 3 via node 2 in the node
6	address fields. All physical and virtual link entries are divided into two
7	fields corresponding to different wavelength channels. In the illustrated
8	example, two channel fields are contained in each link entry.
9	The channel cost of a physical link is a cost associated with a
10	wavelength channel carried by the physical link. Physical link entries 1 and 2
11	are defined for the attributes of the incoming and outgoing links associated
12	with the link switch module 51. In the physical link entry 1 for the incoming
13	link from node 1 to node 2 which uses outgoing interface OG10 and incoming
14	interface IC20, channel cost of 1 is set for both channel #1 and channel #2. In
15	like manner, in the physical link entry 2 for the outgoing link from node 2 to
16	node 3 which uses outgoing interface OG20 and incoming interface IC30,
17	channel cost of 1 is set both for channel #1 and channel #2.
18	Physical link entries 3 and 4 are defined for the attributes of the
19	incoming and outgoing links associated with the channel switch module 52.
20	In the physical link entry 3 for the incoming link from node 1 to node 2 which
21	uses outgoing interface OG11 and incoming interface IC21, channel cost of 1
22	is set for channel #1 and in the physical link entry 4 for the outgoing link
23	from node 2 to node 3 which uses outgoing interface OG21 and incoming
24	interface IC31, channel cost of 1 is set for channel #1.
25	Virtual link entries 1 and 2 are defined for the attributes of virtual links

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1 60 and 64. The channel cost of a virtual link is a total cost of incoming and 2 outgoing physical links and one of the node modules through which the virtual link is established. In the virtual link entry 1 which uses outgoing 3 interface OG10 of node 1 and incoming interface IC30 of node 3, channel cost 4 of 3 (= 1 + 1 + 1) is set for both channels #1 and #2 (λ 1 and λ 2). In the virtual 5 6 link entry 2 which uses outgoing interface OG11 of node 1 and incoming 7 interface IC31 of node 3, channel cost of 4 (= 1 + 2 + 1) is set for channel #1 $(\lambda 3)$. 8 9 The multiplexing of different wavelengths on a single optical link 10 results in the reduction of the number of ports required for group-associated 11 node modules and in the simplification of route computations. Additionally, 12 the provision of the virtual link entries including total channel cost further 13 simplifies route computations. 14 In the previous embodiments, temporary mismatches among the 15 databases of the network nodes 1, 2 and 3 or possible inter-node out-of-16 synchronization may occur. Centralized management of databases and 17 routing control could solve this problem. 18 Fig. 10 shows an embodiment to implement the centralized 19 management with a network management center 65 that performs centralized 20 management of routing control for a plurality of nodes 61, 62 and 63 of 21 identical configuration. A client device 64 is shown connected to the node 62, 22 for example. In Fig. 10, parts corresponding in significance to those in Fig. 1 23 are marked with the same numerals as those in Fig. 1 and the description 24 thereof is omitted for simplicity. Network management center 65 includes a 25 path controller 71, a module state database 72, a link state database 73 and a

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1	terminating unit 74. Path controller 71 collects module and link status
2	information from nodes 61, 62 and 63 via the terminating unit 74 and updates
3	the corresponding databases 72 and 73. To exchange information with the
4	management center 65, each node is provided with a terminating unit 75 to
5	which the client device 64 is also connected. The module and link status
6	information of all network nodes are therefore stored in the databases 72 and
7	73 and maintained by the center 65.
8	When a path setup message is sent from the client device 64 to the
9	node 62, it is retransmitted to the path controller 71. In response, the path
LO	controller 71 determines a minimum cost route using the databases 72 and 73
11	and formulates source and destination control signals based on the
17	determined route and sends the source control signal to the node 62 and the

destination control signal to a node specified in the path setup message.